



FIDO ECDA Algorithm

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Editor:

[Rolf Lindemann](#), [Nok Nok Labs, Inc.](#)

Contributors:

[Jan Camenisch](#), [IBM](#)
[Manu Drijvers](#), [IBM](#)
[Alec Edgington](#), [Trustonic](#)
[Anja Lehmann](#), [IBM](#)
[Rainer Urian](#), [Infineon](#)

The English version of this specification is the only normative version. Non-normative [translations](#) may also be available.

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Abstract

The FIDO Basic Attestation scheme uses attestation "group" keys shared across a set of authenticators with identical characteristics in order to preserve privacy by avoiding the introduction of global correlation handles. If such an attestation key is extracted from one single authenticator, it is possible to create a "fake" authenticator using the same key and hence indistinguishable from the original authenticators by the relying party. Removing trust for registering new authenticators with the related key would affect the entire set of authenticators sharing the same "group" key. Depending on the number of authenticators, this risk might be unacceptable high.

This is especially relevant when the attestation key is primarily protected against malware attacks as opposed to targeted physical attacks.

An alternative approach to "group" keys is the use of individual keys combined with a Privacy-CA [[TPMv1-2-Part1](#)]. Translated to FIDO, this approach would require one Privacy-CA interaction for each Uauth key. This means relatively high load and high availability requirements for the Privacy-CA. Additionally the Privacy-CA aggregates sensitive information (i.e. knowing the relying parties the user interacts with). This might make the Privacy-CA an interesting attack target.

Another alternative is the Direct Anonymous Attestation [[BriCamChe2004-DAA](#)]. Direct Anonymous Attestation is a cryptographic scheme combining privacy with security. It uses the authenticator specific secret once to communicate with a single DAA Issuer and uses the resulting DAA credential in the DAA-Sign protocol with each relying party. The DAA scheme has been adopted by the Trusted Computing Group for TPM v1.2 [[TPMv1-2-Part1](#)].

In this document, we specify the use of an improved DAA scheme based on [[CamDriLeh16-DAA](#)] [[CCDLNU2017-DAA](#)] that uses elliptic curves and bilinear pairings.

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1. Notation

Type names, attribute names and element names are written as `code`.

String literals are enclosed in “”, e.g. “ED256”.

In formulas we use “|” to denote byte wise concatenation operations.

$X = P^x$ denotes scalar multiplication (with scalar x) of a (elliptic) curve point P.

RAND(x) denotes generation of a random number between 0 and x-1.

RAND(G) denotes generation of a random number belonging to Group G.

Specific terminology used in this document is defined in [[FIDOGlossary](#)].

The type `BigNumber` denotes an arbitrary length integer value.

The type `ECPoint` denotes an elliptic curve point with its affine coordinates x and y.

The type `ECPoint2` denotes a point on the sextic twist of a BN elliptic curve over $F(q^2)$. The `ECPoint2` has two affine coordinates each having two components of type `BigNumber`

1.1 Conformance

As well as sections marked as non-normative, all authoring guidelines, diagrams, examples, and notes in this specification are non-normative. Everything else in this specification is normative.

The key words **MUST**, **MUST NOT**, **REQUIRED**, **SHOULD**, **SHOULD NOT**, **RECOMMENDED**, **MAY**, and **OPTIONAL** in this specification are to be interpreted as described in [RFC2119].

2. Overview

This section is non-normative.

FIDO uses the concept of attestation to provide a cryptographic proof of the **authenticator** [FIDOGlossary] model to the relying party. When the **authenticator** is registered to the relying party (RP), it generates a new authentication key pair and includes the public key in the attestation message (also known as key registration data object, **KRD**). When using the ECDAAs algorithm, the **KRD** object is signed using 3.5 ECDAAs-Sign.

For privacy reasons, the authentication key pair is dedicated to one RP (to an application identifier **AppID** [FIDOGlossary] to be more specific). Consequently the attestation method needs to provide the same level of unlinkability. This is the reason why the FIDO ECDAAs Algorithm doesn't use a basename (bsn) often found in other direct anonymous attestation algorithms, e.g. [BriCamChe2004-DAA] or [BFGSW-2011].

The **authenticator** encapsulates all user verification operations and cryptographic functions. An **authenticator** specific module (**ASM**) [FIDOGlossary] is used to provide a standardized communication interface for authenticators. The **authenticator** might be implemented in separate hardware or trusted execution environments. The ASM is assumed to run in the normal operating system (e.g. Android, Windows, ...).

2.1 Scope

This document describes the FIDO ECDAAs attestation algorithm in detail.

2.2 Architecture Overview

ECDAAs attestation defines **global system parameters** and **ECDAAs Issuer specific parameters**. Both parameter sets need to be installed on the host, in the **authenticator** and in the FIDO Server. The ECDAAs method consists of two steps:

- **ECDAAs-Join** between the **authenticator** and the **ECDAAs Issuer** to be performed *before* the first FIDO Registration. The **ECDAAs Issuer** represents the authenticator vendor as it provides the credentials to attest the **authenticator** model.
 - $(n, B, sc, yc) = \text{GetNonceFromECDAAsIssuer}()$
 - $(D=Q, c1, s1) = \text{EcdasJoin1}(X, Y, B, sc, yc, n)$
 - $(A, B, C, D) = \text{EcdasIssuerJoin}(Q, c1, s1)$
 - $\text{EcdasJoin2}(A, C) // \text{store cre}=(A, B, C, D)$
- and the pair of **ECDAAs-Sign** performed by the **authenticator** and **ECDAAs-Verify** performed by the FIDO Server of the relying party as part of the FIDO Registration.
 - Client: $\text{Attestation} = (\text{signature}, \text{KRD}) = \text{EcdasSign}(\text{AppID})$
 - Server: $\text{success} = \text{EcdasVerify}(\text{signature}, \text{KRD}, \text{AppID})$

The technical implementation details of the ECDAAs-Join step are out-of-scope for FIDO. In this document we normatively specify the general algorithm to the extent required for interoperability and we outline examples of some possible implementations for this step.

The ECDAAs-Sign and ECDAAs-Verify steps and the encoding of the related ECDAAs Signature are normatively specified in this document. The generation and encoding of the **KRD** object is defined in other FIDO specifications.

The algorithm and terminology are inspired by [BFGSW-2011]. The algorithm was modified in order to fix security weaknesses (e.g. as mentioned by [ANZ-2013] and [XYZF-2014]). Our algorithm proposes an improved task split for the sign operation while still being compatible to TPMv2 (without fixing the TPMv2 weaknesses in such case).

3. FIDO ECDAAs Attestation

This section is normative.

3.1 Object Encodings

We need to convert `BigNumber` and `ECPoint` objects to byte strings using the following encoding functions:

3.1.1 Encoding `BigNumber` values as byte strings (`BigNumberToB`)

We use the I2OSP algorithm as defined in [RFC3447] for converting big numbers to byte arrays. The bytes from the big endian encoded (non-negative) number `n` will be copied right-aligned into the buffer area `b`. The unused bytes will be set to 0. Negative values will not occur due to the construction of the algorithms.

EXAMPLE 1: Converting `BigNumber n` to byte string `b`

b0 b1 b2 b3 b4 b5 b6 b7

The algorithm implemented in Java looks like this:

EXAMPLE 2: Algorithm for converting BigInteger to byte strings

```

ByteArray BigIntegerToB(
    BigInteger inVal, // IN: number to convert
    int size          // IN: size of the output.
)
{
    ByteArray buffer = new ByteArray(size);
    int oversize = size - inVal.length;
    if (oversize < 0)
        return null;
    for (int i=oversize; i > 0; i--)
        buffer[i] = 0;
    ByteCopy( inVal.bytes, &buffer[oversize], inVal.length);
    return buffer;
}

```

3.1.2 Encoding ECPoint values as byte strings (ECPointToB)

We use the ANSI X9.62 Point-to-Octet-String [ECDSA-ANSI] conversion using the expanded format, i.e. the format where the compression byte (i.e. 0x04 for expanded) is followed by the encoding of the affine x coordinate, followed by the encoding of the affine y coordinate.

EXAMPLE 3: Converting ECPoint P to byte string

```

(x, y) = ECPointGetAffineCoordinates(P)
len = G1.byteLength
byte string = 0x04 | BigIntegerToB(x,len) | BigIntegerToB(y,len)

```

3.1.3 Encoding ECPoint2 values as byte strings (ECPoint2ToB)

The type **ECPoint2** denotes a point on the sextic twist of a BN elliptic curve over $F(q^2)$, see section [4.1 Supported Curves for ECDAAs](#). Each **ECPoint2** is represented by a pair (a, b) of elements of $F(q)$.

The group zero element is always encoded (using the encoding rules as described below) as an element having all components set to zero (i.e. cx.a=0, cx.b=0, cy.a=0, cy.b=0).

We always assume normalized (non-zero) ECPoint2 values (i.e. cz = 1) before encoding them. Non-zero values are encoded using the expanded format (i.e. 0x04 for expanded) followed by the cx followed by the cy value. This leads to the concatenation of 0x04 followed by the first element (cx.a) and second element (cx.b) of the pair of cx followed by the first element (cy.a) and second element (cy.b) of the pair of cy. All individual numbers are padded to the same length (i.e. the maximum byte length of all relevant 4 numbers).

EXAMPLE 4: Converting ECPoint2 P2 to byte string

```

(cx, cy) = ECPointGetAffineCoordinates(P2)
len = G2.byteLength
byte string = 0x04 | BigIntegerToB(cx.a,len) | BigIntegerToB(cx.b,len)
               | BigIntegerToB(cy.a,len) | BigIntegerToB(cy.b,len)

```

3.2 Global ECDAAs System Parameters

1. Groups G_1 , G_2 and G_T , of sufficiently large prime order p
2. Two generators P_1 and P_2 , such that $G_1 = \langle P_1 \rangle$ and $G_2 = \langle P_2 \rangle$
3. A bilinear pairing $e : G_1 \times G_2 \rightarrow G_T$. We propose the use of "ate" pairing (see [BarNae-2006]). For example source code on this topic, see [BNPairings](#).
4. Hash function H with $H : \{0, 1\}^* \rightarrow \mathbb{Z}_p$.
5. (G_1, P_1, p, H) are installed in all authenticator s implementing FIDO ECDAAs attestation.

Definition of G_1 , G_2 , G_T , Pairings, hash function H

See section [4.1 Supported Curves for ECDAAs](#).

3.3 Issuer Specific ECDAAs Parameters

ECDAAs Issuer Parameters consist of the following values:

1. Randomly generated ECDAAs Issuer private key $isk = (x, y)$ with $[x, y] = RAND(p)$.
2. ECDAAs Issuer public key (X, Y) , with $X = P_2^x$ and $Y = P_2^y$.
3. A proof that the ECDAAs Issuer key was correctly computed
 1. BigInteger $r_x = RAND(p)$
 2. BigInteger $r_y = RAND(p)$
 3. ECPoint2 $U_x = P_2^{r_x}$

4. ECPoint2 $U_y = P_2^{r_y}$
 5. BigInteger $c = H(U_x \| U_y \| P_2 \| X \| Y)$
 6. BigInteger $s_x = r_x + c \cdot x \pmod{p}$
 7. BigInteger $s_y = r_y + c \cdot y \pmod{p}$
4. $ipk = X, Y, c, s_x, s_y$

Whenever a party uses ipk for the first time, it must first verify that it was correctly generated:

$$H(P_2^{s_x} \cdot X^{-c} \| P_2^{s_y} \cdot Y^{-c} \| P_2 \| X \| Y) \stackrel{?}{=} c$$

NOTE

$$P_2^{s_x} \cdot X^{-c} = P_2^{r_x + cx} \cdot P_2^{-cx} = P_2^{r_x} = U_x$$

$$P_2^{s_y} \cdot Y^{-c} = P_2^{r_y + cy} \cdot P_2^{-cy} = P_2^{r_y} = U_y$$

The ECDAAs Issuer public key ipk **MUST** be dedicated to a single authenticator model.

We use the element c of ipk as an identifier for the ECDAAs Issuer public key (called **ECDAAs Issuer public key identifier**).

3.4 ECDAAs-Join

NOTE

One ECDAAs-Join operation is required once in the lifetime of an authenticator prior to the first registration of a credential.

In order to use ECDAAs, the authenticator must first receive ECDAAs credentials from an ECDAAs Issuer. This is done by the ECDAAs-Join operation. This operation needs to be performed a single time (before the first credential registration can take place). After the ECDAAs-Join, the authenticator will use the ECDAAs-Sign operation as part of each FIDO Registration. The ECDAAs Issuer is not involved in this step. ECDAAs plays no role in FIDO Authentication / Transaction Confirmation operations.

In order to use ECDAAs, (at least) one ECDAAs Issuer is needed. The approach specified in this document easily scales to multiple ECDAAs Issuer s, e.g. one per authenticator vendor. FIDO lets the authenticator vendor choose any ECDAAs Issuer (similar to his current freedom for selecting any PKI infrastructure/service provider to issuing attestation certificates required for FIDO Basic Attestation).

- All ECDAAs-Join operations (of the related authenticator s) are performed with one of the ECDAAs Issuer entities.
- Each ECDAAs Issuer has a set of public parameters, i.e. ECDAAs public key material. The related Attestation Trust Anchor is contained in the metadata of each authenticator model identified by its AAGUID.

There are two different implementation options relevant for the authenticator vendors (the authenticator vendor can freely choose them):

1. In-Factory ECDAAs-Join
2. Remote ECDAAs-Join and

In the first case, physical proximity is used to locally establish the trust between the ECDAAs Issuer and the authenticator (e.g. using a key provisioning station in a production line). There is no requirement for the ECDAAs Issuer to operate an online web service.

In the second case, some credential is required to remotely establish the trust between the ECDAAs Issuer and the authenticator. As this operation is performed once and only with a single ECDAAs Issuer, privacy is preserved and an authenticator specific credential can and should be used.

Not all ECDAAs authenticator s might be able to add their authenticator model IDs (e.g. AAGUID) to the registration assertion (e.g. TPMs). In all cases, the ECDAAs Issuer will be able to derive the exact the authenticator model from either the credential or the physically proximate authenticator. So the ECDAAs Issuer root key **MUST** be dedicated to a single authenticator model.

3.4.1 ECDAAs-Join Algorithm

This section is normative.

NOTE

If this join is not in-factory, the value Q must be authenticated by the authenticator. Upon receiving this value, the ECDAAs Issuer must verify that this authenticator did not join before.

1. The authenticator asks the ECDAAs Issuer for the B value of the credential.
2. The ECDAAs Issuer chooses a nonce BigInteger $m = \text{RAND}(p)$ such that $x = H(m)$ is on the curve and $yc = \sqrt{(x^3 + b) \bmod q}$.
3. The ECDAAs Issuer computes the B value of the credential as

$$B = (H(m), yc)$$

4. The ECDAAs Issuer sends m and yc to the authenticator.
5. The authenticator verifies that $yc^2 \stackrel{?}{=} H(m)^3 + b$, with b being the b -Value of the related curve.
6. The authenticator chooses and stores the ECDAAs private key $sk = \text{RAND}(p)$
7. The authenticator re-computes $B = (H(m), yc)$
8. The authenticator computes its ECDAAs public key $Q = B^{sk}$
9. The authenticator proves knowledge of sk as follows
 1. BigInteger $r_1 = \text{RAND}(p)$
 2. ECPoint $U_1 = B^{r_1}$
 3. BigInteger $c_2 = H(U_1 || B || Q || m)$
 4. BigInteger $n = \text{RAND}(p)$
 5. BigInteger $c_1 = H(n || c_2)$
 6. BigInteger $s_1 = r_1 + c_1 \cdot sk$
10. The authenticator sends Q, c_1, s_1, n via the ASM to the ECDAAs Issuer
11. The ECDAAs Issuer verifies that the authenticator is "authentic" and that Q was indeed generated by the authenticator. In the case of an in-factory Join, this might be trivial; in the case of a remote Join this typically requires the use of other cryptographic methods. Since ECDAAs-Join is a one-time operation, unlinkability is not a concern for that.
12. The ECDAAs Issuer verifies that $Q \in G_1$ and verifies $H(n || H(B^{s_1} \cdot Q^{-c_1} || B || Q || m)) \stackrel{?}{=} c_1$ (check proof-of-possession of private key).

NOTE

$$B^{s_1} \cdot Q^{-c_1} = B^{r_1 + c_1 sk} \cdot Q^{-c_1} = B^{r_1 + c_1 sk} \cdot B^{-c_1 sk} = B^{r_1} = U_1$$

13. The ECDAAs Issuer creates credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
14. The ECDAAs Issuer sends A, C to the authenticator. The authenticator still knows B and D
15. The authenticator checks that $A, C \in G_1$ and $A \neq 1_{G_1}$
16. The authenticator checks $e(A, Y) \stackrel{?}{=} e(B, P_2)$

NOTE

$$e(A, Y) = e(B^{1/y}, P_2^y) = e(B, P_2^{y/y}) = e(B, P_2);$$

17. and the authenticator checks $e(C, P_2) \stackrel{?}{=} e(A \cdot D, X)$

NOTE

$$e(C, P_2) = e((A \cdot Q)^x, P_2); e(A \cdot D, X) = e(A \cdot Q, P_2^x) = e((A \cdot Q)^x, P_2)$$

18. The authenticator stores credential A, B, C, D

3.4.2 ECDAAs-Join Split between Authenticator and ASM

This section is non-normative.

NOTE

If this join is not in-factory, the value Q must be authenticated by the authenticator. Upon receiving this value, the ECDAAs Issuer must verify that this authenticator did not join before.

1. The ASM asks the ECDAAs Issuer for the B value of the credential.

2. The [ECDAA Issuer](#) chooses a nonce BigInteger $m = \text{RAND}(p)$ such that $x = H(m)$ is on the curve and $yc = \sqrt{((x + b) \bmod q)}$.
3. The [ECDAA Issuer](#) computes the B value of the credential as $B = (H(m), yc)$.
4. The [ECDAA Issuer](#) sends m and yc to the [ASM](#).
5. The [ASM](#) forwards m and yc to the [authenticator](#).
6. The [authenticator](#) verifies that $yc^2 \stackrel{?}{=} H(m)^3 + b$, with b being the b -Value of the related curve.
7. The [authenticator](#) chooses and stores the private key BigInteger $sk = \text{RAND}(p)$.
8. The [authenticator](#) re-computes $B = (H(m), yc)$.
9. The [authenticator](#) computes its ECDAA public key ECPoint $Q = B^{sk}$.
10. The [authenticator](#) proves knowledge of sk as follows
 1. BigInteger $r_1 = \text{RAND}(p)$
 2. ECPoint $U_1 = B^{r_1}$
 3. BigInteger $c_2 = H(U_1 \| B \| Q \| m)$
 4. BigInteger $n = \text{RAND}(p)$
 5. BigInteger $c_1 = H(n \| c_2)$
 6. BigInteger $s_1 = r_1 + c_1 \cdot sk$
11. The [authenticator](#) sends Q, c_1, s_1, n to the [ASM](#), who forwards it to the [ECDAA Issuer](#).
12. The [ECDAA Issuer](#) verifies that the [authenticator](#) is "authentic" and that Q was indeed generated by the [authenticator](#). In the case of an in-factory Join, this might be trivial; in the case of a remote Join this typically requires the use of other cryptographic methods. Since ECDAA-Join is a one-time operation, unlinkability is not a concern for that.
13. The [ECDAA Issuer](#) verifies that $Q \in G_1$ and verifies $H(n \| H(B^{s_1} \cdot Q^{-c_1} \| B \| Q \| m)) \stackrel{?}{=} c_1$.
14. The [ECDAA Issuer](#) creates credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
15. The [ECDAA Issuer](#) sends A, C to the [ASM](#). The ASM remembered B and $D = Q$ from an earlier step.
16. The [ASM](#) checks that $A, B, C, D \in G_1$ and $A \neq 1_{G_1}$.
17. The [ASM](#) checks $e(A, Y) \stackrel{?}{=} e(B, P_2)$
18. and the [ASM](#) checks that $e(C, P_2) \stackrel{?}{=} e(A \cdot D, X)$
19. The [ASM](#) stores A, B, C, D and sends A, C to the [authenticator](#). The [authenticator](#) still knows B and D .
20. The [authenticator](#) stores B, D and ignores further join requests.

NOTE

These values belong to the ECDAA secret key sk . They should persist even in the case of a factory reset.

3.4.3 ECDAA-Join Split between TPM and ASM

This section is non-normative.

NOTE

The Endorsement key credential (EK-C) and TPM2_ActivateCredentials are used for supporting the remote Join.

This description is based on the principles described in [[TPMv2-Part1](#)] section 24 and [Arthur-Challener-2015], page 109 ("Activating a Credential").

1. The [ASM](#) asks the [ECDAA Issuer](#) for the B value of the credential.
2. The [ECDAA Issuer](#) chooses a nonce BigInteger $m = \text{RAND}(p)$ such that $x = H(m)$ is on the curve and $yc = \sqrt{((x^3 + b) \bmod q)}$.
3. The [ECDAA Issuer](#) computes the B value of the credential as $B = (H(m), yc)$.
4. The [ECDAA Issuer](#) sends m and yc to the [ASM](#).
5. The [ASM](#)
 1. instructs the TPM to create a restricted key by calling TPM2_Create, giving the public key template `TPMT_PUBLIC` [TPMv2-Part2] (including the public key P_1 in field `unique`) to the [ASM](#).
 2. re-computes $B = (H(m), yc)$

3. retrieves TPM Endorsement Key Certificate (EK-C) from the TPM
 4. calls $\text{TPM2_Commit}(\text{keyhandle}, P1)$ where keyhandle is the handle of the restricted key generated before (see above), $P1$ is set to $(B.x, B.y)$, and $s2$ and $y2$ are set to $B.x$ and $B.y$ respectively. This call returns K, E , and ctr ; where $K = B^{sk} = Q$, $E = B^{r1}$ is used as U_1 value.
 5. computes $\text{BigInteger } c_2 = H(U_1 \| B \| Q \| m)$
 6. calls $\text{TPM2_Sign}(c_2, \text{ctr})$, returning s_1, n , where $n = \text{RAND}(p)$, and $s_1 = r_1 + c_1 \cdot sk$.
 7. computes $\text{BigInteger } c_1 = H(n \| c_2)$
 8. sends EK-C, TPMT_PUBLIC (including Q in field **unique**), c_1, s_1, n to the [ECDAA Issuer](#).
6. The [ECDAA Issuer](#)
1. verifies EK-C and its certificate chain. As a result the [ECDAA Issuer](#) knows the TPM model related to EK-C.
 2. verifies that this EK-C was not used in a (successful) Join before
 3. Verifies that the **objectAttributes** in TPMT_PUBLIC [\[TPMv2-Part2\]](#) matches the following flags: **fixedTPM** = 1; **fixedParent** = 1; **sensitiveDataOrigin** = 1; **encryptedDuplication** = 0; **restricted** = 1; **decrypt** = 0; **sign** = 1.
 4. examines the public key Q , i.e. it verifies that $Q \in G_1$
 5. checks $H(n \| H(B^{s_1} \cdot Q^{-c_1} \| B \| Q \| m)) \stackrel{?}{=} c_1$
 6. generates the ECDAA credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
 7. generates a **secret** (derived from a **seed**) and wraps the credential A, B, C, D using that **secret**.
 8. encrypts the **seed** using the public key included in EK-C.
 9. uses **seed** and **name** in KDFa (see [\[TPMv2-Part2\]](#) section 24.4) to derive HMAC and **symmetric encryption key**. Wrap the **secret** in **symmetric encryption key** and protect it with the **HMAC key**.

NOTE

The parameter **name** in KDFa is derived from TPMT_PUBLIC , see [\[TPMv2-Part1\]](#), section 16.

10. sends the wrapped object including the credential from previous step to the [ASM](#).
7. The [ASM](#) instructs the TPM (by calling $\text{TPM2_ActivateCredential}$) to
1. decrypt the **seed** using the TPM Endorsement key
 2. compute the **name** (for the ECDAA attestation key)
 3. use the **seed** in KDFa (with **name**) to derive the **HMAC key** and the **symmetric encryption key**.
 4. use the **symmetric encryption key** to unwrap the **secret**.
8. The [ASM](#)
1. unwraps the credential A, B, C, D using the **secret** received from the TPM.
 2. checks that $A, B, C, D \in G_1$ and $A \neq 1_{G_1}$
 3. checks $e(A, Y) \stackrel{?}{=} e(B, P_2)$ and $e(C, P_2) \stackrel{?}{=} e(A \cdot D, X)$
 4. stores A, B, C, D

3.5 ECDAA-Sign

NOTE

One ECDAA-Sign operation is required for the client-side environment whenever a new credential is being registered at a relying party.

3.5.1 ECDAA-Sign Algorithm

This section is normative.

(signature, KRD) = EcdaaSign(String AppID)

Parameters

- p : System parameter prime order of group G_1 (global constant)
- **AppID**: FIDO **AppID** (i.e. https-URL of TrustedFacets object)

Algorithm outline

1. $KRD = \text{BuildAndEncodeKRD}()$; // all traditional Registration tasks are here (e.g. key generation)
2. $\text{BigNumber } l = \text{RAND}(p)$ // use same random number algorithm as for key generation in step 1
3. $\text{ECPoint } R = A^l$;
4. $\text{ECPoint } S = B^l$;
5. $\text{ECPoint } T = C^l$;
6. $\text{ECPoint } W = D^l$;
7. $\text{BigInteger } r = \text{RAND}(p)$ // use same random number algorithm as for key generation in step 1
8. $\text{ECPoint } U = S^r$
9. $\text{BigInteger } c2 = H(U \| S \| W \| \text{AppID} \| H(KRD))$
10. $\text{BigInteger } n = \text{RAND}(p)$ // use same random number algorithm as for key generation in step 1
11. $c = H(n \parallel c2)$
12. $\text{BigInteger } s = r + c \cdot sk \pmod{p}$
13. $\text{signature} = (c, s, R, S, T, W, n)$
14. return (signature, KRD)

3.5.2 ECDAASign Split between Authenticator and ASM

This section is non-normative.

NOTE

This split requires both the authenticator and ASM to be honest to achieve anonymity. Only the authenticator must be trusted for unforgeability. The communication between ASM and authenticator must be secure.

Algorithm outline

1. The ASM randomizes the credential

1. $\text{BigNumber } l = \text{RAND}(p)$

NOTE

All values l **MUST** be unguessable and *unique*, e.g. compute $H(\text{RAND}(p), \text{AppID}, ++\text{seqCnt})$ if in doubt about RAND quality, where seqCnt counts the ECDAASign operations performed by this ASM.

2. $\text{ECPoint } R = A^l$;
3. $\text{ECPoint } S = B^l$;
4. $\text{ECPoint } T = C^l$;
5. $\text{ECPoint } W = D^l$;
2. The ASM sends l, AppID to the authenticator
3. The authenticator performs the following tasks
 1. $KRD = \text{BuildAndEncodeKRD}()$; // all traditional Registration tasks are here (e.g. key generation)
 2. $\text{ECPoint } S' = B^l$
 3. $\text{ECPoint } W' = D^l$
 4. $\text{BigInteger } r = \text{RAND}(p)$ // use same random number algorithm as for key generation in step 3.1
 5. $\text{ECPoint } U = S'^r$
 6. $\text{BigInteger } c2 = H(U \| S' \| W' \| \text{AppID} \| H(KRD))$
 7. $\text{BigInteger } n = \text{RAND}(p)$ // use same random number algorithm as for key generation in step 3.1
 8. $c = H(n \parallel c2)$
 9. $\text{BigInteger } s = r + c \cdot sk \pmod{p}$
 10. Send c, s, KRD, n to the ASM
4. The ASM sets $\text{signature} = (c, s, R, S, T, W, n)$ and outputs (signature, KRD)

3.5.3 ECDAASign Split between TPM and ASM

This section is non-normative.

NOTE

This algorithm is for the special case of a TPMv2 as authenticator. This case requires both the TPM and ASM to be honest for anonymity. Only the TPM must be trusted for unforgeability (see [CCDLNU2017-DAA]).

Algorithm outline

1. The ASM randomizes the credential
 1. BigInteger $l = \text{RAND}(p)$
 2. ECPoint $R = A^l$;
 3. ECPoint $S = B^l$;
 4. ECPoint $T = C^l$;
 5. ECPoint $W = D^l$;
2. The ASM calls TPM2_Commit() with $P1$ set to S and $s2, y2$ empty buffers. The ASM receives the result values $K, L, E = S^r = U$ and ctr. K and L are empty since $s2, y2$ are empty buffers.
3. The ASM calls TPM2_Create to generate the new authentication key pair. The related private key might need to be protected with appropriate access control mechanisms, e.g. see section 8 of [UAFAuthnrCommands].
4. The ASM calls TPM2_Certify() on the newly created key with ctr from the TPM2_Commit and $E = U, S, W, AppID$ as qualifying data. The ASM receives signature value s and related nonce n and attestation block KRD (i.e. TPMS_ATTEST structure in this case).
5. BigInteger $c2 = H(E \| S \| W \| AppID \| H(KRD))$, using KRD as returned by the previous step.
6. The ASM computes: $c = H(n \| c2)$
7. The ASM sets signature = (c, s, R, S, T, W, n) and outputs (signature, KRD)

3.6 ECDA-Verify Operation

This section is normative.

NOTE

One ECDA-Verify operation is required for the FIDO Server as part of each FIDO Registration.

boolean EcdaaVerify(signature, AppID, KRD, ModelName)

Parameters

- p: System parameter prime order of group G_1 (global constant)
- P_2 : System parameter generator of group G_2 (global constant)
- signature: (c, s, R, S, T, W, n)
- AppID: FIDO AppID
- KRD: Attestation Data object as defined in other specifications.
- ModelName: the claimed FIDO authenticator model (i.e. either AAID or AAGUID)

Algorithm outline

1. Based on the claimed ModelName, look up X, Y from trusted source
2. Check that $R, S, T, W \in G_1, R \neq 1_{G_1}$, and $S \neq 1_{G_1}$.
3. $H(n \| H(S^s \cdot W^{-c} \| S \| W \| AppID \| H(KRD))) \stackrel{?}{=} c$; fail if not equal

NOTE

$$B = A^v = P_1^{ly}$$

$$D = Q = B^{sk}$$

$$S = B^l \text{ and } W = D^l$$

$$U = S^r$$

$$\begin{aligned} S^s \cdot W^{-c} &= S^{r+csk} \cdot W^{-c} = U \cdot S^{csk} \cdot W^{-c} \\ &= U \cdot B^{lcsk} \cdot D^{-lc} = U \cdot B^{lcsk} \cdot B^{-lcsk} = U \end{aligned}$$

4. $e(R, Y) = e(S, P_2)$; fail if not equal

NOTE

$$e(R, Y) = e(A^l, P_2^y); e(S, P_2) = e(B^l, P_2) = e(A^{b^l}, P_2)$$

5. $e(T, P_2) \stackrel{?}{=} e(R \cdot W, X)$; fail if not equal

NOTE

$$e(T, P_2) = e(C^l, P_2) = e(A^{xl} \cdot Q^{xlyl_j}, P_2); e(A^l \cdot D^l, X) = e(A^l \cdot Q^{blyl_j}, P_2^x)$$

6. for (all sk' on RogueList) do if $W \stackrel{?}{=} S^{sk'}$ fail;

7. // perform all other processing steps for new credential registration

NOTE

In the case of a TPMv2, i.e. [KRD](#) is a [TPMS_ATTEST](#) object. In this case the verifier must check whether the [TPMS_ATTEST](#) object starts with [TPM_GENERATED](#) magic number and whether its field [objectAttributes](#) contains the flag [fixedTPM](#)=1 (indicating that the key was generated by the TPM).

8. return true;

4. FIDO ECDAAs Object Formats and Algorithm Details

This section is normative.

4.1 Supported Curves for ECDAAs

Definition of G1

G1 is an elliptic curve group $E : y^2 = x^3 + ax + b$ over $F(q)$ with $a = 0$.

Definition of G2

G2 is the p-torsion subgroup of $E'(F_{q^2})$ where E' is a sextic twist of E . With $E' : y'^2 = x'^3 + b'$.

An element of $F(q^2)$ is represented by a pair (a,b) where $a + bX$ is an element of $F(q)[X]/\langle X^2 + 1 \rangle$. We use angle brackets $\langle Y \rangle$ to signify the ideal generated by the enclosed value.

NOTE

In the literature the pair (a,b) is sometimes also written as a complex number $a + b * i$.

Definition of GT

GT is an order-p subgroup of $F_{q^{12}}$.

Pairings

We propose the use of Ate pairings as they are efficient (more efficient than Tate pairings) on Barreto-Naehrig curves [[DevScoDah2007](#)].

Supported BN curves

We use pairing-friendly Barreto-Naehrig [[BarNae-2006](#)] [[ISO15946-5](#)] elliptic curves. The curves [TPM_ECC_BN_P256](#) and [TPM_ECC_BN_P638](#) curves are defined in [[TPMv2-Part4](#)].

BN curves have a Modulus $q = 36 \cdot u^4 + 36 \cdot u^3 + 24 \cdot u^2 + 6 \cdot u + 1$ [[ISO15946-5](#)] and a related order of the group $p = 36 \cdot u^4 + 36 \cdot u^3 + 18 \cdot u^2 + 6 \cdot u + 1$ [[ISO15946-5](#)].

- [TPM_ECC_BN_P256](#) is a curve of form $E(F(q))$, where q is the field modulus [[TPMv2-Part4](#)] [[BarNae-2006](#)]. This curve is identical to the P256 curve defined in [[ISO15946-5](#)] section C.3.5.

NOTE

According to cryptographic research, this curve provides approx. 96 bits of security (see [Crypto and CHES 2016](#)) and hence should NOT be used any longer.

- The values have been generated using $u = -7\ 530\ 851\ 732\ 716\ 300\ 289$.
- Modulus $q = 115\ 792\ 089\ 237\ 314\ 936\ 872\ 688\ 561\ 244\ 471\ 742\ 058\ 375\ 878\ 355\ 761\ 205\ 198\ 700\ 409\ 522\ 629\ 664\ 518\ 163$
- Group order $p = 115\ 792\ 089\ 237\ 314\ 936\ 872\ 688\ 561\ 244\ 471\ 742\ 058\ 035\ 595\ 988\ 840\ 268\ 584\ 488\ 757\ 999\ 429\ 535\ 617\ 037$
- p and q have length of 256 bit each.
- $b = 3$
- $P_{1_256} = (x=1, y=2)$
- $b' = (a=3, b=3)$
- $P_{2_256} = (x, y)$, with
 - $P_{2_256.x} = (a=114\ 909\ 019\ 869\ 825\ 495\ 805\ 094\ 438\ 766\ 505\ 779\ 201\ 460\ 871\ 441\ 403\ 689\ 227\ 802\ 685\ 522\ 624\ 680\ 861\ 435, b=35\ 574\ 363\ 727\ 580\ 634\ 541\ 930\ 638\ 464\ 681\ 913\ 209\ 705\ 880\ 605\ 623\ 913\ 174\ 726\ 536\ 241\ 706\ 071\ 648\ 811)$
 - $P_{2_256.y} = (a=65\ 076\ 021\ 719\ 150\ 302\ 283\ 757\ 931\ 701\ 622\ 350\ 436\ 355\ 986\ 716\ 727\ 896\ 397\ 520\ 706\ 509\ 932\ 529\ 649\ 684, b=113\ 380\ 538\ 053\ 789\ 372\ 416\ 298\ 017\ 450\ 764\ 517\ 685\ 681\ 349\ 483\ 061\ 506\ 360\ 354\ 665\ 554\ 452\ 649\ 749\ 368)$
- **TPM_ECC_BN_P638 [TPMv2-Part4]** uses
 - The values have been generated using $u=365\ 375\ 408\ 992\ 443\ 362\ 629\ 982\ 744\ 420\ 548\ 242\ 302\ 862\ 098\ 433$
 - Modulus $q = 641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 253\ 619\ 096\ 481\ 315\ 470\ 262\ 367\ 432\ 019\ 698\ 642\ 631\ 650\ 152\ 075\ 067\ 922\ 231\ 951\ 354\ 925\ 301\ 839\ 708\ 740\ 457\ 083\ 469\ 793\ 717\ 125\ 223$
 - The related order of the group is $p = 641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 252\ 818\ 101\ 344\ 337\ 098\ 690\ 003\ 906\ 272\ 221\ 387\ 599\ 391\ 201\ 666\ 378\ 807\ 960\ 583\ 525\ 233\ 832\ 645\ 565\ 592\ 955\ 122\ 034\ 352\ 630\ 792\ 289$
 - p and q have length of 638 bit each.
 - $b = 257$
 - $P_{1_638} = (x=641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 253\ 619\ 096\ 481\ 315\ 470\ 262\ 367\ 432\ 019\ 698\ 642\ 631\ 650\ 152\ 075\ 067\ 922\ 231\ 951\ 354\ 925\ 301\ 839\ 708\ 740\ 457\ 083\ 469\ 793\ 717\ 125\ 222, y=16)$
 - $b' = (a=771, b=1542)$
 - $P_{2_638} = (x, y)$, with
 - $P_{2_638.x} = (a=192\ 492\ 098\ 325\ 059\ 629\ 927\ 844\ 609\ 092\ 536\ 807\ 849\ 769\ 208\ 589\ 403\ 233\ 289\ 748\ 474\ 758\ 010\ 838\ 876\ 457\ 636\ 072\ 173\ 883\ 771\ 602\ 089\ 605\ 233\ 264\ 992\ 910\ 618\ 494\ 201\ 909\ 695\ 576\ 234\ 119\ 413\ 319\ 303\ 931\ 909\ 848\ 663\ 554\ 062\ 144\ 113\ 485\ 982\ 076\ 866\ 968\ 711\ 247, b=166\ 614\ 418\ 891\ 499\ 184\ 781\ 285\ 132\ 766\ 747\ 495\ 170\ 152\ 701\ 259\ 472\ 324\ 679\ 873\ 541\ 478\ 330\ 301\ 406\ 623\ 174\ 002\ 502\ 345\ 930\ 325\ 474\ 988\ 134\ 317\ 071\ 869\ 554\ 535\ 111\ 092\ 924\ 719\ 466\ 650\ 228\ 182\ 095\ 841\ 246\ 668\ 361\ 451\ 788\ 368\ 418\ 036\ 777\ 197\ 454\ 618\ 413\ 255)$
 - $P_{2_638.y} = (a=622\ 964\ 952\ 935\ 200\ 827\ 531\ 506\ 751\ 874\ 167\ 806\ 262\ 407\ 152\ 244\ 280\ 323\ 674\ 626\ 687\ 789\ 202\ 660\ 794\ 092\ 633\ 841\ 098\ 984\ 322\ 671\ 973\ 226\ 667\ 873\ 503\ 889\ 270\ 602\ 870\ 064\ 426\ 165\ 592\ 237\ 410\ 681\ 318\ 519\ 893\ 784\ 898\ 821\ 343\ 051\ 339\ 820\ 566\ 224\ 981\ 344\ 169\ 470, b=514\ 285\ 963\ 827\ 225\ 043\ 076\ 463\ 721\ 426\ 569\ 583\ 576\ 029\ 220\ 880\ 138\ 564\ 906\ 219\ 230\ 942\ 887\ 639\ 456\ 599\ 654\ 554\ 743\ 732\ 087\ 558\ 187\ 149\ 207\ 036\ 952\ 474\ 092\ 411\ 405\ 629\ 612\ 957\ 921\ 369\ 286\ 372\ 038\ 525\ 830\ 610\ 755\ 207\ 588\ 843\ 864\ 366\ 759\ 521\ 090\ 861\ 911\ 494)$
- **ECC_BN_DSD_P256 [DevScoDah2007]** section 3 uses
 - The values have been generated using $u=6\ 917\ 529\ 027\ 641\ 089\ 837$
 - Modulus $q = 82434016654300679721217353503190038836571781811386228921167322412819029493183$
 - The related order of the group is $p = 82434016654300679721217353503190038836284668564296686430114510052556401373769$
 - p and q have length of 256 bit each.
 - $b = 3$
 - $P_{1_DSD_P256} = (1, 2)$
 - $b' = (a=3, b=6)$
 - $P_{2_DSD_P256} = (x, y)$, with
 - $P_{2_DSD_P256.x} = (a=73\ 481\ 346\ 555\ 305\ 118\ 071\ 940\ 904\ 527\ 347\ 990\ 526\ 214\ 212\ 698\ 180\ 576\ 973\ 201\ 374\ 397\ 013\ 567\ 073\ 039, b=28\ 955\ 468\ 426\ 222\ 256\ 383\ 171\ 634\ 927\ 293\ 329\ 392\ 145\ 263\ 879\ 318\ 611\ 908\ 127\ 165\ 887\ 947\ 997\ 417\ 463)$
 - $P_{2_DSD_P256.y} = (a=3\ 632\ 491\ 054\ 685\ 712\ 358\ 616\ 318\ 558\ 909\ 408\ 435\ 559\ 591\ 759\ 282\ 597\ 787\ 781\ 393\ 534\ 962\ 445\ 630\ 353, b=60\ 960\ 585\ 579\ 560\ 783\ 681\ 258\ 978\ 162\ 498\ 088\ 639\ 544\ 584\ 959\ 644\ 221\ 094\ 447\ 372\ 720\ 880\ 177\ 666\ 763)$
- **ECC_BN_ISO_P512 [ISO15946-5]** section C.3.7 uses
 - The values have been generated using $u=138\ 919\ 694\ 570\ 470\ 098\ 040\ 331\ 481\ 282\ 401\ 523\ 727$
 - Modulus $q = 13\ 407\ 807\ 929\ 942\ 597\ 099\ 574\ 024\ 998\ 205\ 830\ 437\ 246\ 153\ 344\ 875\ 111\ 580\ 494\ 527\ 427\ 714\ 590\ 099\ 881\ 795\ 845\ 981\ 157\ 516\ 604\ 994\ 291\ 639\ 750\ 834\ 285\ 779\ 043\ 186\ 149\ 750\ 164\ 319\ 950\ 153\ 126\ 044\ 364\ 566\ 323$
 - The related order of the group is $p = 13\ 407\ 807\ 929\ 942\ 597\ 099\ 574\ 024\ 998\ 205\ 830\ 437\ 246\ 153\ 344\ 875\ 111\ 580\ 494\ 527\ 427\ 714\ 590\ 099\ 881\ 680\ 053\ 891\ 920\ 200\ 409\ 570\ 720\ 654\ 742\ 146\ 445\ 677\ 939\ 306\ 408\ 461\ 754\ 626\ 647\ 833\ 262\ 056\ 300\ 743\ 149$
 - p and q have length of 512 bit each.
 - $b = 3$
 - $P_{1_ISO_P512} = (x=1, y=2)$
 - $b' = (a=3, b=3)$
 - $P_{2_ISO_P512} = (x, y)$, with

- $P_2_ISO_P512.x = (a=3\ 094\ 648\ 157\ 539\ 090\ 131\ 026\ 477\ 120\ 117\ 259\ 896\ 222\ 920\ 557\ 994\ 037\ 039\ 545\ 437\ 079\ 729\ 804\ 516\ 315\ 481\ 514\ 566\ 156\ 984\ 245\ 473\ 190\ 248\ 967\ 907\ 724\ 153\ 072\ 490\ 467\ 902\ 779\ 495\ 072\ 074\ 156\ 718\ 085\ 785\ 269, b=3\ 776\ 690\ 234\ 788\ 102\ 103\ 015\ 760\ 376\ 468\ 067\ 863\ 580\ 475\ 949\ 014\ 286\ 077\ 855\ 600\ 384\ 033\ 870\ 546\ 339\ 773\ 119\ 295\ 555\ 161\ 718\ 985\ 244\ 561\ 452\ 474\ 412\ 673\ 836\ 012\ 873\ 126\ 926\ 524\ 076\ 966\ 265\ 127\ 900\ 471\ 529)$
- $P_2_ISO_P512.y = (a=7\ 593\ 872\ 605\ 334\ 070\ 150\ 001\ 723\ 245\ 210\ 278\ 735\ 800\ 573\ 263\ 881\ 411\ 015\ 285\ 406\ 372\ 548\ 542\ 328\ 752\ 430\ 917\ 597\ 485\ 450\ 360\ 707\ 892\ 769\ 159\ 214\ 115\ 916\ 255\ 816\ 324\ 924\ 295\ 339\ 525\ 686\ 777\ 569\ 132\ 644\ 242, b=9\ 131\ 995\ 053\ 349\ 122\ 285\ 871\ 305\ 684\ 665\ 648\ 028\ 094\ 505\ 015\ 281\ 268\ 488\ 257\ 987\ 110\ 193\ 875\ 868\ 585\ 868\ 792\ 041\ 571\ 666\ 587\ 093\ 146\ 239\ 570\ 057\ 934\ 816\ 183\ 220\ 992\ 460\ 187\ 617\ 700\ 670\ 514\ 736\ 173\ 834\ 408)$

NOTE

Spaces are used inside numbers to improve readability.

Hash Algorithm H

Depending on the curve, we use $H(x) = \text{SHA256}(x) \bmod p$ or $H(x) = \text{SHA512}(x) \bmod p$ as hash algorithm $H: \{0, 1\}^* \rightarrow Z_p$.

The argument of the hash function must always be converted to a byte string using the appropriate encoding function specific in section 3.1 [Object Encodings](#), e.g. according to section 3.1.3 [Encoding ECPoint2 values as byte strings \(ECPoint2ToB\)](#) in the case of `ECPoint2` points.

NOTE

We don't use [IEEE P1363.3](#) section 6.1.1 IHF1-SHA with security parameter t (e.g. $t=128$ or 256) as it is more complex and not supported by TPMv2.

4.2 ECDAAs Algorithm Names

We define the following JWS-style algorithm names (see [\[RFC7515\]](#)):

ED256

`TPM_ECC_BN_P256` curve, using SHA256 as hash algorithm H .

ED256-2

`ECC_BN_DSD_P256` curve, using SHA256 as hash algorithm H .

ED512

`ECC_BN_ISOP512` curve, using SHA512 as hash algorithm H .

ED638

`TPM_ECC_BN_P638` curve, using SHA512 as hash algorithm H .

4.3 ecdaaSignature object

The fields c and s both have length N . The fields R , S , T , W have equal length $(2*N+1)$ each).

In the case of `BN_P256` curve (with key length $N=32$ bytes), the fields R , S , T , W have length $2*32+1=65$ bytes. The fields c and s have length $N=32$ each.

The `ecdaaSignature` object is a binary object generated as the concatenation of the binary fields in the order described below (total length of 356 bytes for 256bit curves):

Value	Length (in Bytes)	Description
UINT8[] ECDAASignature_c	N	The c value, $c = H(n \parallel c2)$ as returned by <code>EcdaaSign</code> encoded as byte string according to <code>BigNumberToB</code> . Where • $c2 = H(U \parallel S \parallel W \parallel \text{KRD} \parallel \text{AppID})$ • $U = S^r$, with $r = \text{RAND}(p)$ computed by the signer. • KRD is the the entire to-be-signed object (e.g. <code>TAG_UAFV1_KRD</code> in the case of FIDO UAF). • $S = B^l$, with $l = \text{RAND}(p)$ computed by the signer and $B = A^v$ computed in the ECDAAs-Join
UINT8[] ECDAASignature_s	N	The s value, $s = r + c * sk \pmod p$, as returned by <code>EcdaaSign</code> encoded as byte string according to <code>BigNumberToB</code> . Where • $r = \text{RAND}(p)$, computed by the signer at FIDO registration (see 3.5.2 ECDAAs-Sign Split between Authenticator and ASM) • p is the group order of $G1$ • sk: is the <code>authenticator</code>'s attestation secret key, see above
UINT8[]		

ECDAASignature_n	N	The Nonce value n, as returned by Ecdasign encoded as byte string according to BigIntegerToB.
UINT8[] ECDAASignature_R	2*N+1	$R = A^l$; computed by the ASM or the authenticator at FIDO registration; encoded as byte string according to ECPublicToB. Where 1 = RAND(p), i.e. random number $0 \leq l \leq p$. Computed by the ASM or the authenticator at FIDO registration. - And where $R = A^l$ denotes the scalar multiplication (of scalar l) of a curve point A. - Where A has been provided by the ECDAASigner as part of ECDAASign: $A = B^{1/y}$, see 3.4.1 ECDAASign Algorithm. - Where p is a system value, injected into the authenticator and y is part of the ECDAASigner private key $isk=(x,y)$.
UINT8[] ECDAASignature_S	2*N+1	$S = B^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPublicToB. Where B has been provided by the ECDAASigner on Join: $B = (H(m), yC)$, see 3.4.1 ECDAASign Algorithm .
UINT8[] ECDAASignature_T	2*N+1	$T = C^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPublicToB. Where $C = (A \cdot Q)^x$, provided by the ECDAASigner on Join - x is a components of the ECDAASigner private key, $isk=(x,y)$. - Q is the authenticator public key
UINT8[] ECDAASignature_W	2*N+1	$W = D^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPublicToB. Where $D = Q$ is computed by the ECDAASigner at Join (see 3.4.1 ECDAASign Algorithm).

5. Considerations

This section is non-normative.

A detailed security analysis of this algorithm can be found in [[CamDriLeh16-DAA](#)].

5.1 Algorithms and Key Sizes

The proposed algorithms and key sizes are chosen such that compatibility to TPMv2 is possible.

5.2 Indicating the Authenticator Model

Some [authenticator](#) s (e.g. TPMv2) do not have the ability to include their model (i.e. vendor ID and model name) in attested messages (i.e. the to-be-signed part of the registration assertion). The TPM's endorsement key certificate typically contains that information directly or at least it allows the model to be derived from the endorsement key certificate.

In FIDO, the relying party expects the ability to cryptographically verify the [authenticator](#) model.

We require the [ECDAASigner](#) s public key ($ipk=(X,Y,c,sx,sy)$) to be dedicated to one single [authenticator](#) model (e.g. as identified by AAID or AAGUID).

5.3 Revocation

If the private ECDAASign key sk of an [authenticator](#) has been leaked, it can be revoked by adding its value to a RogueList.

The ECDAASign-Verifier (i.e. FIDO Server) check for such revocations. See section [3.6 ECDAASign-Verify Operation](#).

The [ECDAASigner](#) is expected to check revocation by other means:

1. if ECDAASign is done in-factory, it is assumed that produced devices are known to be uncompromised (at time of production).
2. if a remote ECDAASign is performed, the (remote) [ECDAASigner](#) already must use a different method to remotely authenticate the [authenticator](#) (e.g. using some endorsement key). We expect the [ECDAASigner](#) to perform a revocation check based on that information. This is even more flexible as it does not require access to the [authenticator](#) ECDAASign private key sk .

5.4 Pairing Algorithm

The pairing algorithm [e](#) needs to be used by the [ASM](#) as part of the Join process and by the verifier (i.e. FIDO relying party) as part of the verification (i.e. FIDO registration) process.

The result of such a pairing operation is only compared to the result of another pairing operation computed by the same entity. As a consequence, it doesn't matter whether the [ASM](#) and the verifier use the exact same pairings or not (as long as they both use valid pairings).

5.5 Performance

For performance reasons the calculation of $\text{Sig2}=(R, S, T, W)$ may be performed by the [ASM](#) running on the FIDO user device (as opposed to inside the authenticator). See section [3.5.2 ECDAASign Split between Authenticator and ASM](#).

The cryptographic computations to be performed inside the authenticator are limited to G1. The ECDAASigner has to perform two G2 point multiplications for computing the public key. The Verifier (i.e. FIDO relying party) has to perform G1 operations and two pairing operations.

5.6 Binary Concatenation

We use a simple byte-wise concatenation function for the different parameters, i.e. $H(a,b) = H(a \parallel b)$.

This approach is as secure as the underlying hash algorithm since the authenticator controls the length of the (fixed-length) values (e.g. U, S, W). The AppID is provided externally and has unverified structure and length. However, it is only followed by a fixed length entry - the (system defined) hash of KRD. As a consequence, no parts of the AppID would ever be confused with the fixed length value.

5.7 IANA Considerations

This specification registers the algorithm names "ED256", "ED512", and "ED638" defined in section [4. FIDO ECDAASign Object Formats and Algorithm Details](#) with the IANA JSON Web Algorithms registry as defined in section "Cryptographic Algorithms for Digital Signatures and MACs" in [\[RFC7518\]](#).

Algorithm Name	"ED256"
Algorithm Description	FIDO ECDAASign algorithm based on TPM_ECC_BN_P256 [TPMv2-Part4] curve using SHA256 hash algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance , Contact Us
Specification Documents	Sections 3. FIDO ECDAASign Attestation and 4. FIDO ECDAASign Object Formats and Algorithm Details of [FIDOECDAASignAlgorithm] .
Algorithm Analysis Document(s)	[CamDriLeh16-DAA]

Algorithm Name	"ED512"
Algorithm Description	ECDAASign algorithm based on ECC_BN_ISOP512 [ISO15946-5] curve using SHA512 algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance , Contact Us
Specification Documents	Sections 3. FIDO ECDAASign Attestation and 4. FIDO ECDAASign Object Formats and Algorithm Details of [FIDOECDAASignAlgorithm] .
Algorithm Analysis Document(s)	[CamDriLeh16-DAA]

Algorithm Name	"ED638"
Algorithm Description	ECDAASign algorithm based on TPM_ECC_BN_P638 [TPMv2-Part4] curve using SHA512 algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance , Contact Us
Specification Documents	Sections 3. FIDO ECDAASign Attestation and 4. FIDO ECDAASign Object Formats and Algorithm Details of [FIDOECDAASignAlgorithm] .
Algorithm Analysis Document(s)	[CamDriLeh16-DAA]

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